Databases

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When you write a program all your data disappears after the program ends

Unless we save it somewhere

SQL Databases are a sensible choice for where to save your data

- Highly optimized storage of tabular data
- ► Fast and well understood query language
- ► Fault tolerant protocols

So what is a database?

Super fancy spreadsheet

- Each database will contain tables that store data
- Data in tables can be *queried* using a language called SQL
- Data in tables can be *joined* with data in other tables to answer questions

Designing them so you don't tie yourself in knots is tricky!

So why not just use Spreadsheets?

See Matt Parker's excellent Stand-up Maths video: UK Government loses data because of Excel mistake.



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https://youtu.be/zUp8pkoeMss

Different types of database

Traditionally, the database would reside on a separate machine

Space is expensive!

If you wanted to use the database you had to connect to it But nowadays space is cheap

Local per app databases very common

If you need remote data access:

▶ Use a server-style database like MariaDB or MySQL

Otherwise use a file-style database:

▶ ...just use SQLite.

Should I use a database?

Am I being paid to store/process this data?

- ▶ Yes? Use a database.
- ▶ No? Use a spreadsheet (or a database)

Does the data need to be accessed remotely?

- Yes? Use a server-style database (MySQL/MariaDB)
- ▶ No? Use a file-style database (SQLite)

Am I just playing with data or is my data tiny (gigabytes in size)?

▶ Yes? Use a database or plain text data storage (i.e. CSV).

Is my data really big (petabytes in size)?

▶ Yes? Use a NoSQL database (beyond scope of this course)

Does my data contain recursive data structures (i.e. lists of lists of arbitrary length)

Yes? Use Prolog or Datalog. (or abuse a database ; -))

Databases let us store data in tables!

- But how do you structure your data in a table?
- And can we draw pretty doodles based on them?

Relational modelling

Proviso!

Relational modelling is a tool for thinking about how to decompose relationships between things into tables.

People get fussy about the syntax

Please don't!

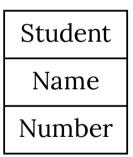
I'll try and show you various syntaxes you may encounter, but its just a tool

- Do whatever works for you
- So long as its clear it doesn't matter
- ▶ The diagrams are for doodling ideas *not* final implementation

Things are nouns!

Here is a student! Students have a name and a number!

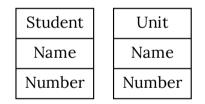
- ▶ The student is the *entity*.
- ▶ The name and number are the *attributes*.



More things are nouns!

Here is a unit! Units also! have a name and a number!

- ▶ The unit is the *entity*.
- ▶ The name and number are the *attributes*.



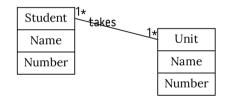
Don't worry about names

There may be many examples of different *values* that could be examples of units and students... but don't worry about that.

| | Student | Unit | |
|--------|------------------|---------------------|-----------|
| Name | Patrick McGoohan | Name Software Tools | |
| Number | 6 | Number | COMS10012 |

Nouns can be related!

One student may take many units; and units may have many students



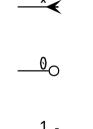
Alternative notation

Some people prefer a graphical notation for entity relationships called *crow's* foot

► I prefer to write it explicitly

Don't get too hung up on notation!

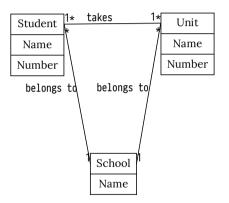
- And use a key if you're ever asked in an exam
- ▶ The point is to let *you* doodle notes
- Do whatever makes sense to you or the people you work with



Schools are a thing!

There are things called schools:

- Schools have names
- Each unit belongs to *exactly* one school
- Each student belongs to *exactly* one school Each school can have students and units its responsible for
 - But could also be empty!



What should I call a student?

Obviously their name would be *polite*... ...but what will happen if we were to open a class on *Gallifrey*?



All 12!

This would rapidly get too confusing for computers!

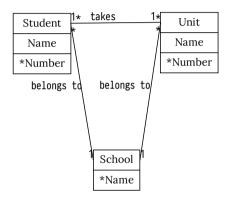
► (But not for people)

A key for an *entity* is the set of attributes needed to uniquely refer to it.

- A candidate key is a minimal set of attributes needed to uniquely refer to it.
- The primary key for an entity is the key we use.

If a key contains multiple attributes its called a *composite key*.

If a key is a meaningless ID column you added just for the sake of having a key its called a *surrogate key*.



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When we want to turn it into tables

Every entity becomes a table

Each table has a primary key

Every *edge* becomes a table

- Contents of these tables are the *primary* keys of the two items being linked
- Attribute that refers to another key is called a *foreign key*

School Membership

| Student | School |
|---------|----------------------------|
| 6970 | School of Computer Science |

School Units

UnitSchoolCOMS10012School of Computer Science

Student

| Name | Number |
|----------------|--------|
| Joseph Hallett | 6970 |

Unit

| Name | Number |
|----------------|-----------|
| Software Tools | COMS10012 |

School

Name School of Computer Science

Class Register

Student Unit 6970 COMS10012

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Conclusions

1. Doodle entity relationship diagrams to sketch out database designs

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- **2**. Convert to databases by making everything a table
- **3**. Don't get hung up on notation

We've got a database for storing data...

► It'd be nice to be able to acutally use it and make queries! For that we need SQL:

Structured Query Language

Query language for asking questions about databases from 1974

Standardized in 1986 in the US and 1987 everywhere else

► Still the dominant language for queries today

Not a general purpose programming language

- ▶ Not Turing complete
- Weird English-like syntax

Standardized?

You would be so lucky!

- ▶ In theory, yes
- ▶ In practice, absolutely not

Every database engine has small differences ...

Some have quite big ones too!

Lots have differences in performance

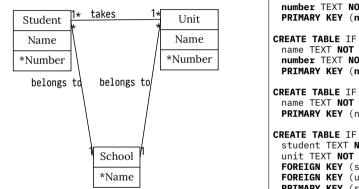
SQLite is good with strings, most others prefer numbers

Managing these differences used to be an entire degree/job in its own right!

Now we just manage databases badly!

I'll try and stick to SQLite's syntax...

CREATE TABLE



CREATE TABLE IF NOT EXISTS student (name TEXT NOT NULL, number TEXT NOT NULL, PRIMARY KEY (number)); CREATE TABLE IF NOT EXISTS unit (

name TEXT NOT NULL, number TEXT NOT NULL, PRIMARY KEY (number));

CREATE TABLE IF NOT EXISTS school (
 name TEXT NOT NULL,
 PRIMARY KEY (name));

CREATE TABLE IF NOT EXISTS class_register (
 student TEXT NOT NULL,
 unit TEXT NOT NULL,
 FOREIGN KEY (student) REFERENCES student(number),
 FOREIGN KEY (unit) REFERENCES unit(name),
 PRIMARY KEY (student, unit));

Lets build it in SQL



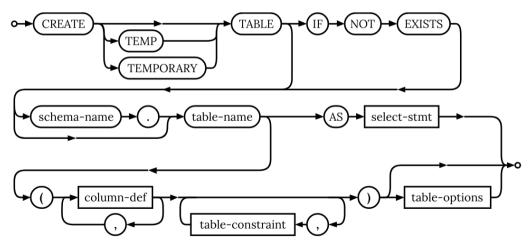
What about if we want to delete them?

| DROP | TABLE | IF | EXISTS | class_register; |
|------|-------|----|--------|-----------------|
| DROP | TABLE | IF | EXISTS | student; |
| DROP | TABLE | IF | EXISTS | unit; |
| DROP | TABLE | IF | EXISTS | school; |
| | | | | |

Syntax, syntax, syntax

If you go on the SQLite documentation page...

- ...you can find syntax diagrams for all of SQL!
- https://www.sqlite.org/lang_createtable.html



Types

When creating the fields in our database we made them all of type TEXT...

► What other types exist?

INTEGER whole numbers REAL lossy decimals BLOB binary data (images/audio/files...) VARCHAR(10) a string of 10 characters TEXT any old text BOOLEAN True or false DATE Today DATETIME Today at 2pm

But really types

Databases sometimes simplify these types

SQLite makes the following tweaks...

INTEGER whole numbers **REAL** lossy decimals **BLOB** binary data (images/audio/files...) VARCHAR(10) actually TEXT **TEXT** any old text **BOOLEAN** actually INTEGER **DATE** actually TEXT DATETIME actually TEXT (others may exist... read the manual!)

Table constraints

In the earlier examples we marked some columns as NOT $\,$ NULL $\,$

- Others as PRIMARY KEY and others as FOREIGN KEY...
- ...what other constraints have we got

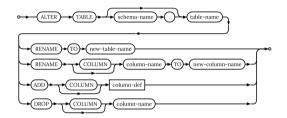
...but SQLite won't actually enforce any of these types or constraints unless you ask it to :- (

 Check out the STRICT keyword when creating the table. NOT NULL can't be NULL
 UNIQUE can't be the same as another row
 CHECK arbitrary checking (including it conforms to a regular expression)
 PRIMARY KEY unique, not NULL and (potentially) autogenerated
 FOREIGN KEY (IGNORED BY MARIADB) other key must exist

Can I add constraints later?

Yes with the ALTER TABLE statement

- But often easiest just to save the table somewhere else
- Drop the table
- Reimport it



What about if we want to add data to a table?

INSERT INTO unit(name, number)
VALUES ("Software_Tools", "COMS100012");

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So far

We've introduced how to:

- ► CREATE TABLE
- ► DROP TABLE
- ► INSERT INTO

Next step: querying data!

I'm going to use a database from an old iTunes library for demo purposes

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Chinook database

SELECT

Basic command for selecting rows from a table is ${\sf SELECT}$

| SELECT * FROM album LIMIT 5; | | | LIMIT 5; | | | |
|---------------------------------|---------------------------------------|----------|----------|----------|-------------------|--|
| | | | | ArtistId | Name | |
| | | | | 1 | AC/DC | |
| AlbumId | Title | ArtistId | | 2 | Accept | |
| 1 | For Those About To Rock We Salute You | 1 | _ | 3 | Aerosmith | |
| 2 | Balls to the Wall | 2 | | 4 | Alanis Morissette | |
| 3 | Restless and Wild | 2 | | 5 | Alice In Chains | |
| 4 | Let There Be Rock | 1 | | | | |
| 5 | Big Ones | 3 | | | | |

SELECT + EDOM artist

JOIN

Ideally we'd like those two tables combined into one...

SELECT * FROM album JOIN artist ON album.artistid = artist.artistid LIMIT 5;

| AlbumId | Title | ArtistId | ArtistId | Name |
|---------|---------------------------------------|----------|----------|-----------|
| 1 | For Those About To Rock We Salute You | 1 | 1 | AC/DC |
| 2 | Balls to the Wall | 2 | 2 | Accept |
| 3 | Restless and Wild | 2 | 2 | Accept |
| 4 | Let There Be Rock | 1 | 1 | AC/DC |
| 5 | Big Ones | 3 | 3 | Aerosmith |

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Reducing the columns...

Clearly there are too many columns here... lets only select the ones we need

SELECT album.title, artist.name FROM album JOIN artist ON album.artistid = artist.artistid LIMIT 5;

| Title | Name |
|---------------------------------------|-----------|
| For Those About To Rock We Salute You | AC/DC |
| Balls to the Wall | Accept |
| Restless and Wild | Accept |
| Let There Be Rock | AC/DC |
| Big Ones | Aerosmith |

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Renaming columns

Title and Name aren't particularly meaningful without context

Lets name them something sensible

| album | artist |
|---------------------------------------|-----------|
| For Those About To Rock We Salute You | AC/DC |
| Balls to the Wall | Accept |
| Restless and Wild | Accept |
| Let There Be Rock | AC/DC |
| Big Ones | Aerosmith |

I'm feeling rocky

I want to listen to something a bit rocky...

▶ Lets filter all the albums to the ones that have Rock in the title

| album | artist |
|---------------------------------------|-------------|
| For Those About To Rock We Salute You | AC/DC |
| Let There Be Rock | AC/DC |
| Deep Purple In Rock | Deep Purple |
| Rock In Rio [CD1] | Iron Maiden |
| Rock In Rio [CD2] | Iron Maiden |

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Who rocks?

So who has put out an album with Rock in it?

```
SELECT artist.name AS artist
FROM album
JOIN artist
ON album.artistid = artist.artistid
WHERE album.title LIKE '%Rock%'
LIMIT 5;
```

| artist |
|-------------|
| AC/DC |
| AC/DC |
| Deep Purple |
| Iron Maiden |
| Iron Maiden |

| SELECT DISTINCT artist.name AS artist |
|--|
| FROM album |
| JOIN artist |
| ON album.artistid = artist.artistid |
| WHERE album.title LIKE '%Rock%' |
| LIMIT 5; |

| artist |
|--------------------|
| AC/DC |
| Deep Purple |
| Iron Maiden |
| The Cult |
| The Rolling Stones |

How many rock albums has each artist put out?

Lets group by artist and count the albums!

SELECT artist.name AS artist, COUNT(album.title) as albums FROM album JOIN artist ON album.artistid = artist.artistid WHERE album.title LIKE '%Rock%' GROUP BY artist LIMIT 5;

| artist | albums |
|--------------------|--------|
| AC/DC | 2 |
| Deep Purple | 1 |
| Iron Maiden | 2 |
| The Cult | 1 |
| The Rolling Stones | 1 |

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Really we want this list ordered...

Lets group by artist and count the albums...

And order it by album count!

```
SELECT artist.name AS artist,
        COUNT(album.title) as albums
FROM album
JOIN artist
ON album.artistid = artist.artistid
WHERE album.title LIKE '%Rock%'
GROUP BY artist
ORDER BY albums DESC
LIMIT 5;
```

| artist | albums |
|--------------------|--------|
| Iron Maiden | 2 |
| AC/DC | 2 |
| The Rolling Stones | 1 |
| The Cult | 1 |
| Deep Purple | 1 |

Basics of SQL

So thats the basics of SQL!

- ▶ You can do a *bunch* more things with SQL SELECT statements...
- ...you can pick them up as you write queries.
- ...most SQL engines have a bunch more counting and query functions too

Go read the documentation!

Database theory!

- ▶ So far we've discussed how to doodle database designs...
- We've discussed how to create tables in SQL

How do we design tables that are easy to use?

Lets start with our records database again...

We could store our data as follows:

| Artist | Albums |
|---------------|--|
| The Beatles | Yellow Submarine, White Album, Rubber Soul |
| Milk Can | Make It Sweet |
| Dresden Dolls | Yes Virginia, No Virginia, The Dresden Dolls |

Please, no.

This is a terrible idea

- Yes we have one big table which seems neater
- But its much harder to do anything actually with

For example:

- How many albums does each artist have?
- ▶ Change all of Prince's albums after 1993 to being by a Love Symbol

How many artists have an album with the same name?

Normal forms

Normal forms prevent this sort of insanity

- ▶ Using them requires discipline, and remembering rules...
- But is worth it for your sanity in the short to medium term

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Each column shall contain one (and only one) value

Each row says describes *multiple* albums per artist...

| Artist | Albums |
|---------------|--|
| The Beatles | Yellow Submarine, White Album, Rubber Soul |
| Milk Can | Make It Sweet |
| Dresden Dolls | Yes Virginia, No Virginia, The Dresden Dolls |

First Normal Form

Lets fix that...

| A | rtist | Album |
|---|---------------|-------------------|
| Г | The Beatles | Yellow Submarine |
| Г | The Beatles | White Album |
| Г | The Beatles | Rubber Soul |
| Ν | ⁄lilk Can | Make It Sweet |
| Γ | Dresden Dolls | Yes Virginia |
| Γ | Dresden Dolls | No Virginia |
| Γ | Dresden Dolls | The Dresden Dolls |

Lets add some more data to our table

| Artist | Album | Year | Prime Minister |
|---------------|-------------------|------|----------------|
| The Beatles | Yellow Submarine | 1969 | Harold Wilson |
| The Beatles | White Album | 1968 | Harold Wilson |
| The Beatles | Rubber Soul | 1965 | Harold Wilson |
| Milk Can | Make It Sweet | 1999 | Tony Blair |
| Dresden Dolls | Yes Virginia | 2006 | Tony Blair |
| Dresden Dolls | No Virginia | 2008 | Gordon Brown |
| Dresden Dolls | The Dresden Dolls | 2003 | Tony Blair |

Second Normal Form

Every non-key attributue is fully dependent on the key

In this case the key is Artist, Album

► And arguably *year* too if you're gonna pull a Taylor Swift and rerelease all your albums... Is *Prime Minister* dependent on the key?

▶ No. Put it in a different table.

Now it looks like

| Artist | Album | Year |
|---------------|-------------------|------|
| The Beatles | Yellow Submarine | 1969 |
| The Beatles | White Album | 1968 |
| The Beatles | Rubber Soul | 1965 |
| Milk Can | Make It Sweet | 1999 |
| Dresden Dolls | Yes Virginia | 2006 |
| Dresden Dolls | No Virginia | 2008 |
| Dresden Dolls | The Dresden Dolls | 2003 |

| Year | Prime Minister |
|------|----------------|
| 1969 | Harold Wilson |
| 1968 | Harold Wilson |
| 1965 | Harold Wilson |
| 1999 | Tony Blair |
| 2006 | Tony Blair |
| 2008 | Gordon Brown |
| 2003 | Tony Blair |

Third Normal Form

Every non-key attribute must provide a fact about the key, the whole key and nothing but the key; so help me Codd.

Lets add some extra information to our table of Prime Ministers...

| Year | Prime Minister | Birthday |
|------|----------------|------------|
| 1969 | Harold Wilson | 1916-03-11 |
| 1968 | Harold Wilson | 1916-03-11 |
| 1965 | Harold Wilson | 1916-03-11 |
| 1999 | Tony Blair | 1953-05-06 |
| 2003 | Tony Blair | 1953-05-06 |
| 2006 | Tony Blair | 1953-05-06 |
| 2008 | Gordon Brown | 1951-02-20 |

Our key is (Year, Prime Minister); Birthday depends on Prime Minister.

- So every non-key depends on the key...
- ► So 2NF

But not 3NF as Birthday doesn't tell you a fact about the whole key... just the Prime Minister.

So split it up!

| Year | Prime Minister |
|------|----------------|
| 1969 | Harold Wilson |
| 1968 | Harold Wilson |
| 1965 | Harold Wilson |
| 1999 | Tony Blair |
| 2003 | Tony Blair |
| 2006 | Tony Blair |
| 2008 | Gordon Brown |
| | |

| Prime Minister | Birthday |
|----------------|------------|
| Harold Wilson | 1916-03-11 |
| Tony Blair | 1953-05-06 |
| Gordon Brown | 1951-02-20 |

Why is this better?

- Now if we need to alter the birthday of a PM (or any other fact about that key)...
- ...then we only need to alter it in one place.

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Other normal forms...

Boyce-Codd Normal Form

A slightly stronger form of 3NF...

Sometimes called 3.5th Normal Form

Every possible candidate key for a table is also in 3NF.

Split a 3NF table into tables with single candidate keys to get 3.5NF.

4th Normal Form

If multiple attributes in a table depend on the same key,

- Then those attributes should be dependent too
- Otherwise split them into separate tables...

5th Normal Form

It's in 4th normal form and you can't split it into more separate tables.

This is all getting a bit mathsy...

You can look up formal definitions for each of the normal forms

(and you should)

But so long as you keep things as separate as possible, you'll usually hit at least 3NF by accident.

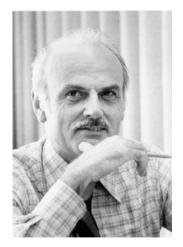
- ...and practically speaking your probably good then
- ▶ Getting it to 5NF does make things more flexible in the long run...
- ▶ But a 3.5NF database is often *good enough*.

Ultimately design is subjective (somewhat).

...but mathematical proof of flexibility is good right?

Codd's Law

Every non-key attribute must provide a fact about the key, the whole key and nothing but the key; so help me Codd.



Lets get back to SQL

So far we've introduced basic SQL

- How to create tables
- How to add and delete data
- ► How to run basic queries

Lets go further!

► More features, more function

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- ► Other joins
- NULL

There is a special value in SQL to represent missing data: NULL.

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- But they're pretty much always a bad idea
- ▶ The logic for comparing them is pretty whacky

NULL = NULL?

Lets say we have a database with the following table:

| Person | Fruit |
|---------|-------|
| Joseph | Lime |
| Matt | Apple |
| Manolis | |

Lets find everyone who we know what their favourite fruit is!

| | SELECT | * | FROM | fruit | WHERE | fruit | <> | NULL; | |
|--|--------|---|------|-------|-------|-------|----|-------|--|
|--|--------|---|------|-------|-------|-------|----|-------|--|

Err..., lets try the opposite?

SELECT * FROM fruit WHERE fruit = NULL;

Err what?

SELECT * FROM fruit WHERE fruit LIKE '%';

| Person | Fruit |
|--------|-------|
| Joseph | Lime |
| Matt | Apple |

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So...

SELECT * FROM fruit WHERE fruit NOT LIKE '%';

NULL is weird...

Because NULL means attribute missing...

► The results of comparing with it are just plain stupid somewhat unexpected The simple solution is to declare everything as NOT NULL

► And use a higher normal form (5NF) then you'll find they almost entirely disappear Otherwise you have to memorise a bunch of stupid special comparators

SELECT * FROM fruit WHERE fruit IS NULL;

Person Fruit Manolis

SELECT * FROM fruit WHERE fruit IS NOT NULL;

Person Fruit Joseph Lime Matt Apple

Tricky joins

Clearly testing for equality when NULL is problematic.

▶ So what happens when you want to join two tables together with NULL's in them

| Person | Fruit | | Fruit | Dish |
|---------|-------|---|--------|---------------|
| Joseph | Lime | - | Apple | Apple crumble |
| Matt | Apple | | Banana | Banana split |
| Manolis | | | Cherry | |
| | | | Lime | Daiquiri |

So what might make a nice dish for each of your lecturers?

 (A NATURAL JOIN is like a regular JOIN but assumes same named columns ought to be equal).

| Person | Fruit | Dish |
|--------|-------|---------------|
| Joseph | Lime | Daiquiri |
| Matt | Apple | Apple crumble |

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But what about poor Manolis? How do we get him to appear in our table?

LEFT and RIGHT JOIN

When doing our previous JOIN we wanted only rows that matched...

► Technically called an INNER JOIN...

Sometimes we're okay with the database sticking NULL in if we want to keep columns where a join *can't* be made...

SELECT person, fruit.fruit, dish
FROM fruit
LEFT JOIN recipes
ON fruit.fruit = recipes.fruit;

| Person | Fruit | Dish |
|---------|-------|---------------|
| Joseph | Lime | Daiquiri |
| Matt | Apple | Apple crumble |
| Manolis | | |

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RIGHT JOIN A RIGHT JOIN is like a left join but the other way round...

SELECT fruit.fruit, dish, person
FROM fruit
RIGHT JOIN recipes
ON fruit.fruit = recipes.fruit;

| Fruit | Dish | Person |
|-------|---------------|--------|
| Lime | Daiquiri | Joseph |
| Apple | Apple crumble | Matt |
| | Banana split | |

Where has the Banana gone?!

SELECT recipes.fruit, dish, person
FROM fruit
RIGHT JOIN recipes
ON fruit.fruit = recipes.fruit;

| Fruit | Dish | Person |
|--------|---------------|--------|
| Lime | Daiquiri | Joseph |
| Apple | Apple crumble | Matt |
| Banana | Banana split | |
| Cherry | - | |

(Or just NATURAL JOIN and it'll usually take care of it...)

SELECT fruit, dish, person FROM fruit RIGHT NATURAL JOIN recipes;

| Fruit | Dish | Person | | |
|--------|------------------------|-------------|-----|-----|
| Lime | Daiquiri | Joseph | | |
| Apple | Apple crumble | Matt | | |
| Banana | Banana split | | | |
| Cherry | < □ ¹ < @ > | 《 문 》 《 문 》 | - E | SQ(|

One more JOIN!

What if we want to do a LEFT and a RIGHT JOIN at the same time?

SELECT *
FROM fruit
FULL OUTER NATURAL JOIN recipes;

| Person | Fruit | Dish |
|---------|------------------|---------------|
| Joseph | Lime | Daiquiri |
| Matt | Apple | Apple crumble |
| Manolis | | |
| | Banana Cherry | Banana split |

What about statistic functions?

In the last lecture we introduced COUNT as a way of counting how many things exist?

▶ How may different fruits are in the outer joined table?

SELECT *
FROM fruit
FULL OUTER NATURAL JOIN recipes;

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| SELEC | ст сои | NT(fruit) NATURAL |) | |
|-------|--------|----------------------|------|---------|
| FROM | fruit | | | |
| FULL | OUTER | NATURAL | JOIN | recipes |

| Person | Fruit | Dish |
|---------|--------|---------------|
| Joseph | Lime | Daiquiri |
| Matt | Apple | Apple crumble |
| Manolis | | |
| | Banana | Banana split |
| | Cherry | |

COUNT(fruit)

...So it looks like COUNT ignores NULL

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Other statistics...

Lets rank fruits!

| Fruit | Stars |
|--------|-------|
| Apple | 0 |
| Banana | 4 |
| Cherry | |
| Lime | 5 |

SELECT AVG(stars) AS Average FROM ranking;

Average 3.0

SELECT SUM(stars)/COUNT(fruit) AS Average
FROM ranking;

Average 2

Remember computers are awful

- Multiply count by 1.0 to "fix"?
- Also number of stars is ordinal data so the mean shouldn't be used anyway...

What about standard deviation?

The standard deviation is how far something deviates on average from the mean.

```
SELECT SQRT(AVG(Deviation)) AS STDDEV
FROM (
   SELECT Fruit, Stars, Mean,
      (Stars-Mean)*(Stars-Mean) AS Deviation
FROM ranking JOIN (
      SELECT AVG(stars) AS Mean
      FROM ranking
   )
   WHERE stars IS NOT NULL
);
```

STDDEV

2.16024689946929

You can nest queries inside one another (subqueries!)

- ▶ This is a recipe for making your SQL *slow*
- Maybe just use SQL for data retrieval and leave complex stats to statistical programming languges?

So thats SQL!

Tips for using it?

- Don't overcomplicate things!
- ▶ Normal forms make things simpler!

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► Avoid NULL like the plague

In the real world we rarely want to access a database in its own right

► Rather it is used within a programming language as part of a program Different languages have different APIs for different databases...

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...but Java has the JDBC for almost all of them

- Library is in java.sql and javax.sql packages
- ▶ Wraps all of a databases functionality into something that looks a lot like Oracle SQL.

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Supports prepared statements (you want to use these)

```
import java.sql.*;
try (final Connection conn = DriverManager.getConnection("jdbc:sqlite:database.db")) {
    conn.createStatement()
        .executeQuery("CREATE_TABLE_users(username_TEXT_PRIMARY_KEY, _password_TEXT)");
} catch (final SQLException err) {
    System.out.println(err);
}
```

java.sql.SQLException: No suitable driver found for jdbc:sqlite:database.db

Lets add some suitable users...

```
import java.sql.*;
import java.util.*;
final var users = new HashMap<String, String>();
users.put("Joseph", "password");
users.put("Matt", "password1");
users.put("Manolis", "12345");
try (final Connection conn = DriverManager.getConnection("jdbc:sqlite:database.db")) {
   conn.createStatement().executeUpdate("DELETE,FROM,users"):
   final var statement = conn.prepareStatement("INSERT, INTO, users, VALUES(?,..?)");
   for (final var user : users.keySet()) {
      statement.setString(1, user);
      statement.setString(2, users.get(user));
      statement.executeUpdate();
catch (final SOLException err) {
  System.out.println(err);
3
```

And list them back out...

| Matt | password1 |
|---------|-----------|
| Joseph | password |
| Manolis | 12345 |

Why not this...

When adding all the users we used a PreparedStatement to add all the users.

```
final var statement = conn.prepareStatement("INSERT_INTO_USERS_VALUES(?,..?)");
for (final var user : users.keySet()) {
    statement.setString(1, user);
    statement.setString(2, users.get(user));
    statement.executeUpdate();
}
```

Wouldn't this be easier?

```
for (final var user : users.keySet())
    conn.createStatement()
        .executeUpdate("INSERT_INTO_users_"+"VALUES_('"+user+"',__'"+users.get(user)+"')");
```

SQL Injection

This leads to a horrible vulnerability called an injection attack

- ▶ You can do something similar with shellscript too ;)
- Search for Shellshock vulnerability if you're interested...

What a prepared statement does is ensure that the things you add are what you say they are Suppose you do the something similar for the login code:

SELECT username FROM users
WHERE username = "Joseph"
AND password = "password";

username

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Joseph

Suppose the username and password are taken from a website login form...

- What happens if I try and login with a password of:
- " OR 1 OR password = "heheh

Bad things

With a prepared statement:

SELECT username FROM users
WHERE username = "Joseph"
AND password = """_OR_1_OR_password_=_""heheh";

Without a prepared statement:

SELECT username FROM users
WHERE username = "Joseph"
AND password = "" OR 1 OR password = "heheh";

Matt Joseph Manolis

ALWAYS USE PREPARED STATEMENTS

The compiler will even spew warnings and errors about this nowadays...

Transactions

Another cool thing that JDBC makes easy are *transactions*... Suppose you want to do a bunch of additions and updates to a database...

▶ What happens if something goes wrong *in the middle*?

You could go and manually roll back all the new data you added and changes you made ...

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- Sounds tedious
- Lets automate it!

Transaction workflow

- **1**. Start a new transaction
- **2**. Do your work
- **3**. Commit to it when done
- 4. Rollback if an error occurs

And in Java please?

```
import java.sql.*;
import java.util.*:
try (final Connection conn = DriverManager.getConnection("jdbc:sglite:database.db")) {
   conn.setAutoCommit(false):
   final var save = conn.setSavepoint();
  trv ₹
     conn.createStatement() .executeQuery("INSERT_INTO_users_VALUES_('Alice', 'pa55w0rd')");
     conn.createStatement() .executeOuerv("INSERT, INTO, users, VALUES, ('Bob', 'Pa55w0Rd7')");
     if (true) throw new Exception("Whoops!");
      conn.createStatement() .executeQuery("INSERT_INTO_users_VALUES_('Eve',_'backd00r')");
     conn.commit():
   catch (final Exception err) {
     conn.rollback(save);
   finallv 
     conn.setAutoCommit(true):
   ۲
catch (final SOLException err) {
  System.out.println(err);
```

Now if we query users...

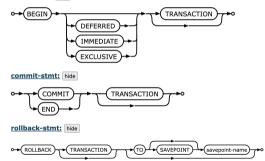
SELECT * FROM users;

| username | password |
|------------------|-----------------------|
| Matt | password1 |
| Joseph | password |
| Manolis | 12345 |
| | |
| | |
| username | password |
| username Matt | password password1 |
| | |

Our table *remains* unaltered... the whole transaction was *rolled back*.

(Oh, and BTW SQLite also can do transactions in SQL)

begin-stmt: hide



Conclusions

JDBC lets you access SQL from Java

- Make sure you load the right driver
- Catch SQLExceptions
- Use prepared statements and transactions to prevent errors
- And an ORM like Hibernate if you like.

IMPORTANT NOTE

Please don't actually implement password storage like we did in the lecture...

- ▶ Go speak to someone in the cyber or crypto groups first...
- Or read NIST 800-63 first

I will write papers about you if you do ; -)

Joseph Hallett, Nikhil Patnaik, Benjamin Shreeve and Awais Rashid. "Do this! Do that!, And nothing will happen" Do specifications lead to securely stored passwords? 2021 IEEE/ACM 43rd International Conference on Software Engineering (ICSE). 2021.